



# Towards Deployment-Efficient Reinforcement Learning: Lower Bound and Optimality

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• An Abstraction of Learning Process of Online RL





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# Background



Learn a good policy with K as small as possible, while N can be large.



• An Abstraction of Learning Process of Online RL



We expect to learn a good policy with K as small as possible, while N can be large.

# A New Theoretical Formulation for DE-RL

#### • [Definition] Deployment Complexity

- An algorithm has *deployment complexity K*, if for arbitrary MDP,  $\epsilon$ ,  $\delta > 0$ , the algorithm:
  - (1) return  $\epsilon$ -optimal policy w.p.  $1 \delta$  after K policy switching,
  - (2) collect *N* trajectories in each deployment, with the constraint that *N* is **polynominal level in standard parameters**.

To avoid trivial case

- Overall Goal of Deployment-Efficient RL (DE-RL)
  - Q1: Lower bound of the deployment complexity?
  - Q2: Any algorithms matching lower bound?

## Related Work

#### Empirical Literature

• [1] propose the conception of "Deployment-Efficient RL", while a clear theoretical formulation is still an open problem.

#### • Theoretical Literature (low-switching cost RL)

- Tabular MDP [2] and Linear MDP [3]
- Relies on "adaptive policy switching strategy", unrealistic when policy deployment is timeconsuming.
- Consider to reduce # of policy switching while achieving near-optimal regret;
- Only study deploying deterministic policy each time;

[1] Matsushima et. al., 2020, Deployment-Efficient Reinforcement Learning via Model-Based Offline Optimization
[2] Bai et. al., 2020, Provably efficient q-learning with low switching cost
[3] Gao et. al., 2021, A provably efficient algorithm for linear markov decision process with low switching cost

Result in sub-optimal deployment complexity

### Linear MDP as a Concrete Example

#### • Episodic (Time-dependent) Linear MDP

- Horizon Length: *H*
- State-action feature:  $\phi(s_h, a_h) \in \mathbb{R}^d$ , with  $||\phi(s_h, a_h)|| \le 1$
- Reward is linear:  $r_h(s, a) = \phi^T(s_h, a_h)\theta_h$
- Transition is linear:  $P_h(s_{h+1}|s_h, a_h) = \phi^T(s_h, a_h)\mu_h(s_{h+1})$

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#### • [Definition] Deployment Complexity in Linear MDP

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#### Separately consider two settings

- **Case 1**: Can only deploy deterministic policy;
  - Setting in previous low-switching cost setting
- **Case 2**: Allow to deploy arbitrary policy (deterministic/stochastic/non-Markovian);
  - Largely omitted in previous literatures, but strictly lower deployment complexity.

## Lower Bound of Deployment Complexity

- **Case 1**: Only allow to deploy deterministic policy;
  - Lower bound:  $\Omega(dH)$
  - Intuition: a (*d*, *H*)-linear MDP has *dH* different "dimensions to explore", while each deterministic policy can only explore one of them.
- **Case 2**: Allow to deploy arbitrary policy;
  - Lower bound:  $\Omega(H/\log N)$
  - Intuition: the agent can only push exploration forward for  $O(\log N)$  time steps (layers).

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- Corollary for time-homogeneous linear MDP:
  - $\Omega(d)$  for Case 1 (deterministic policy setting)
  - Ω(min{*d*, *H*}/log *N*) for Case 2 (arbitrary policy setting)

# Algorithm with Near-Optimal Deployment Complexity

- High-level idea
  - Explore each layer (time step) per O(d) or O(1) deployments







# Algorithm with Near-Optimal Deployment Complexity

- Case 1: Only allow to deploy deterministic policy;
  - Algorithm: LSVI-UCB [Jin et. al., 2020] + Layer-by-layer
  - A novel elliptical potential lemma (Lem. 4.2)
  - Guarantees:
    - Deployment complexity *K*: *O*(*dH*);
    - (Asmptotically) # of Trajs  $N: O(\frac{H^4 d^3}{\epsilon^2} \log^2 \frac{H d}{\delta \epsilon})$
  - Remarks:
    - Layer-by-layer strategy is not necessary for deployment efficiency in case 1, but some additional benefits (see Appx. C.4).
    - Similar analysis can be extended to reward-free setting

# Algorithm with Near-Optimal Deployment Complexity

- **Case 2**: Allow to deploy arbitrary policy;
  - Algorithm: a new batch exploration algorithm
    - Explore in a layer-by-layer manner
    - For each layer
      - Use a novel covariance matrix estimation method to evaluate the exploration ability of given policies;
      - Plus a **bonus-term driven** method to find a set of deterministic policies **cover all the dimensions**.
  - Require an additional reachability assumption  $v_{min}$ .
  - Guarantees:
    - Deployment complexity  $K: \Theta(H)$ ;
    - (Asmptotically) # of Trajs N: Poly $(d, H, \frac{1}{\epsilon}, \log \frac{1}{\delta}, \frac{1}{v_{\min}})$
  - Remarks:
    - Naturally a reward-free exploration
    - Open problem: is it possible to remove dependence on  $\nu_{min}$

# Thanks!